# COMP3153/9153 Homework 1

Temporal Logic, CTL Model Checking, Büchi Automata

**Due: Fri 20th March 2020, 10am** Submission guidelines are given at the end of this document.

# **Exercise 1 (LTL – Part I)**

### (20 Marks)

Assume the set of atomic propositions is  $\{a, b, c\}$ .

**Question 1** For each of the following set of paths  $S_i$ ,  $1 \le i \le 5$ , give an LTL formula,  $\varphi_i$ , that best defines  $S_i$  (and justify your answer):

- 1. S<sub>1</sub>: paths that contain at least two a;
- 2. S<sub>2</sub>: paths that do not contain a state satisfying b and  $\neg a$ ;
- 3. S<sub>3</sub>: paths that contain at least one c which is not immediately followed neither by a nor b;
- 4. S<sub>4</sub>: paths that eventually do not contain a a.
- 5.  $S_5$ : paths that satisfy: if containing a finite number of a then they do not contain an infinite number of c.

Note: you can use the operators **F** and **G**, as well as all other LTL-operations.

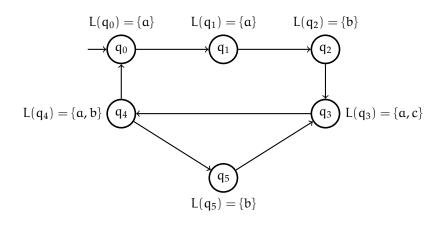


Figure 1: Automaton A

**Question 2** Given the labeled automaton A of Fig. 1, with labeling-function L. For each of the following LTL formulas, is there at least one run/trace in A that satisfies the formula? If you answer "Yes" give a witness, and if you answer "No" justify your answer.

1. <b>X X</b> b;	4. <b>FG</b> (( <b>X</b> ¬b) ∨ b);
2. <b>F</b> c;	5. <b>G</b> (a <b>UNTIL</b> $(\neg a \land b)$ ).
3. <b>G F G</b> (b ∨ a);	

**Question 3** Among the LTL formulas (1–5) of Question 2, which are the ones satisfied by automaton A? (Explain your answer).

# Exercise 2 (CTL – Part I) (10 Marks)

A system is composed of two tanks, A and B, and a pump Pump. The states of the system are labelled by atomic propositions from the set

{pumpOn, pumpOff, tankAEmpty, TankAFull, TankBEmpty, tankBFull}

**Question 4** Consider the following informal requirements and CTL formulas that should formally define the corresponding requirements:

- 1. "It is always the case that if the pump is off, it will be on again in the future." CTL formula: **AG** (pumpOff ⇒ **AF** pumpOn)
- "The pump is always off if tank A is empty or tank B is full".
  CTL formula: AG AF (pumpOff ⇒ (tankAEmpty ∨ tankBFull))
- 3. *"It is always possible to reach a state where tank B is full".* CTL formula: **AG EF** tankBFull

One of the previous CTL formulas does not match the informal requirement. Which one? (Explain your answer).

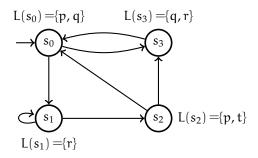


Figure 2: Automaton B

**Question 5** What would be the correct CTL formula corresponding to the informal requirement?

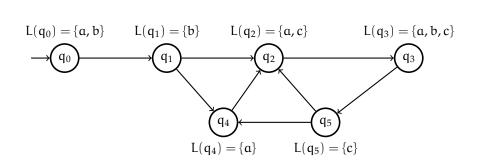
## Exercise 3 (CTL — Part II) (20 Marks)

**Question 6** Consider the automaton B of Figure 2. Check whether,  $B, s_0 \models \phi$  and  $B, s_2 \models \phi$  for the following CTL formulas  $\phi$ :

- 1. **AF** q
- 2. **AG**(**EF**  $(p \lor r)$ )
- 3. **EX**(**EX** r)
- $4. \ \textbf{AG}(\textbf{AF} \ q)$

**Question 7** Which of the following pairs of CTL formulas are equivalent. For those that are not equivalent, find a model of one of the pair which is not a model of the other.

- 1. **EF**  $\phi$  and **EG**  $\phi$
- 2.  $(\textbf{EF} \varphi) \lor (\textbf{EF} \psi)$  and  $\textbf{EF} (\varphi \lor \psi)$
- 3.  $(AF \varphi) \lor (AF \psi)$  and  $AF (\varphi \lor \psi)$
- 4. **AF**  $\neg \phi$  and  $\neg$ **EG**  $\phi$
- 5. true and  $(\mathbf{AG} \ \varphi) \Rightarrow (\mathbf{EG} \ \varphi)$
- 6. true and  $(\textbf{EG} \varphi) \Rightarrow (\textbf{AG} \varphi)$



(30 Marks)

Figure 3: Automaton C

**Question 8** Let  $\varphi$  be the following CTL formula:

**Exercise 4 (CTL Model Checking)** 

$$\mathsf{AG}\left(b \Rightarrow \mathsf{AX}\neg \mathsf{E}\big((a \lor c) \mathsf{UNTIL} b\big)\right).$$

- 1. Rewrite the formula  $\varphi$  into equivalent formula  $\varphi'$  without **AG**, **AX** and  $\Rightarrow$ . We consider that the *or* operator,  $\lor$ , is allowed in a CTL formula.
- 2. Give the parse tree of the rewritten formula  $\varphi'$ .
- 3. Manually run the CTL explicit-state marking algorithm from Week 2 on  $\varphi'$  i.e., compute the result of Mark( $C, \varphi'$ ), to determine whether  $C \models \varphi'$  (and explain your answer).

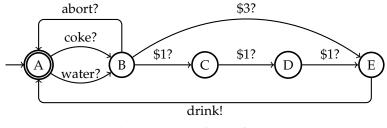
### Exercise 5 (Büchi Automata)

Figures 4(a) and 4(b) model a vending machine and a customer, respectively. The automata are *Büchi automata*.

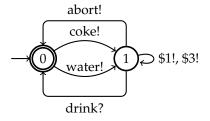
The two automata synchronise on matching send/receive actions: a!/a? results in a common action of the components named a.

**Question 9** Briefly describe the behaviour of the automata in English.

**Question 10** Build the synchronised product *Vending Machine2* × *Customer2*.



(a) Automaton Vending Machine2



(b) Automaton Customer2

Figure 4: Automata for Exercise 5

#### **Submission Guidelines**

- Due time: Fri 20th March 2020, 10am. No late submission allowed.
- Submit one PDF file (hw1.pdf) using the CSE give system by typing the command give cs3153 hw1 hw1.pdf on a CSE terminal. Alternatively use the online submission page.
- It is highly recommended that you use LATEX to prepare your document. A guide is provided on the course website.